On partitioning graphs into connected subgraphs

Julien Bensmail

Université Nice-Sophia-Antipolis, France

COATI seminar
April 4th, 2017

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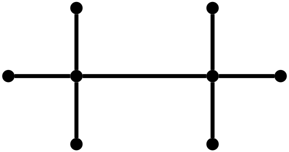

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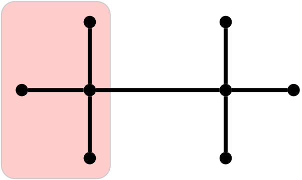

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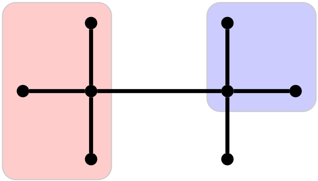

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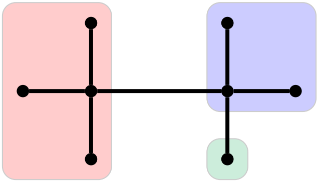

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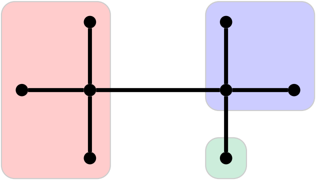

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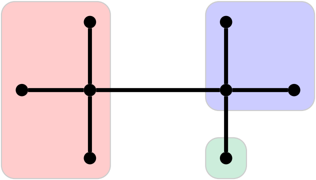

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$(V_1, \ldots, V_p)$ is a **realization** of $(n_1, \ldots, n_p)$ in G.

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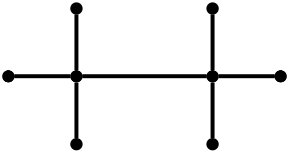

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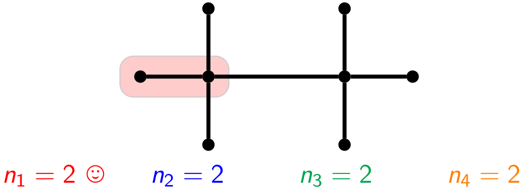

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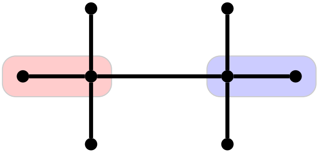

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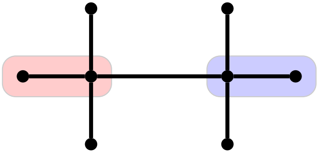

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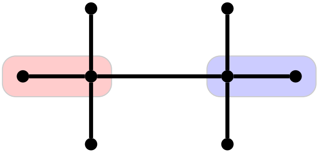

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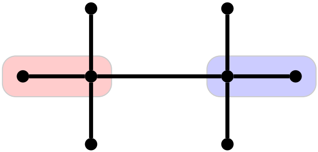

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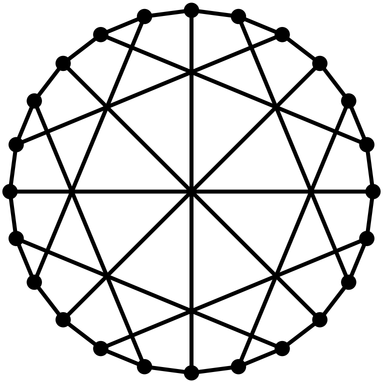

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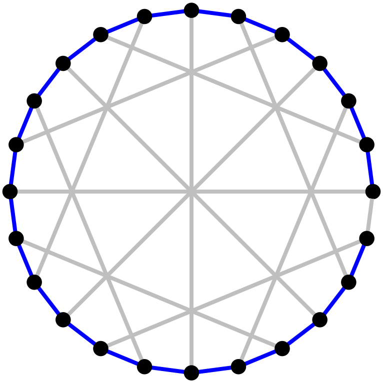

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So far, considered aspects of AP graphs include:

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Many open questions...

Algorithmic aspects

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REALIZATION
Input: A graph G, and a partition $\pi := (n_1, ..., n_p)$ of $|V(G)|$.
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REALIZATION is NP-complete, even under many restrictions:

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   - when $|\mathrm{sp}(\pi)| = 1$ (i.e. $\pi = (k, ..., k)$ for $k \geq 3$) [Dyer, Frieze, 1985];
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   - when G is a tree with $\Delta(G) = 3$ [Barth, Fournier, 2006];
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So, what about the problem $\mathrm{AP} = \{$Graph G: is G AP?$\}$?

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$\Rightarrow$ Generally yield checking algorithms.

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What for other classes of graphs?

(e.g. general trees, 3-connected near-triangulations, etc.)

Structural aspects

On AP trees

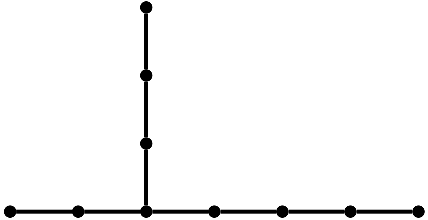
AP trees are rather understood:

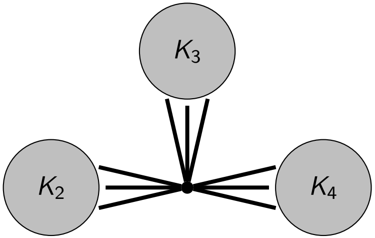
Theorem [Barth, Fournier, Ravaux, 2009]

- AP trees have $\Delta \leq 4$;
- degrees at least 3 are located on a same path;
- degree-4 vertices are adjacent to a leaf.

Rephrased differently...

Obtained by considering surgraphs that are "easier" w.r.t. the AP property:





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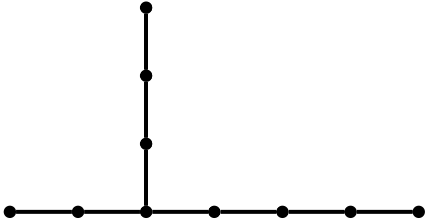

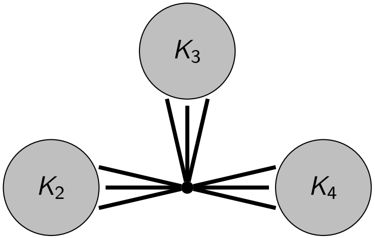

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So, actually:

Corollary [Barth, Fournier, Ravaux, 2009]

Removing a cut-vertex from an AP graph results in at most 4 components.

Generalization

Via the same technique:

Theorem [Baudon, Foucaud, Przybyło, Woźniak, 2014]

For any $k \geq 2$, removing a k-cutset from an AP graph:

- may result in arbitrarily many components,
- whose orders grow exponentially.

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Let G be a graph with order n. If for every two non-adjacent vertices u and v of G we have $d(u) + d(v) \geq n - 1$, then G is traceable.

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Let G be a graph with order $n \geq 8$. If $\alpha(G) \leq \lceil n/2 \rceil$ and for every two non-adjacent vertices u and v of G we have $d(u) + d(v) \geq n - 3$, then G is AP.

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Brandt claimed a generalization to triples of independent vertices.

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Theorem [Kalinowski, Pilśniak, Schiermeyer, Woźniak, 2016]

Let G be a connected graph with order $n \geq 22$. If $|E(G)| > \binom{n-4}{2} + 12$, then G is AP.

Further directions

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- For $k \geq 3$, k-connected k-regular are not all traceable...
- ... what for AP graphs? [Diwan, 2003]
- Pick your favourite result on traceability. Does it weaken to AP graphs?
  - Closure?
  - Sets of forbidden patterns?
  - etc.

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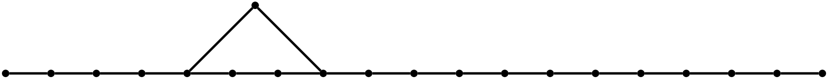

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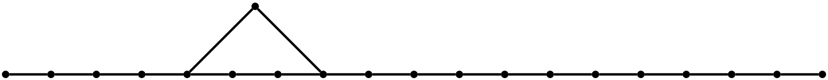

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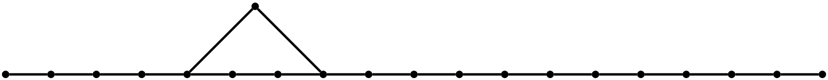

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Properties of minimal AP graphs?

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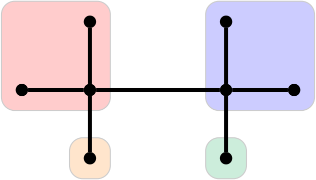
Questions:

- Denser families?
- Generalization of the Δ property.
- Clique number?
- Families with connectivity $k \geq 2$?
- etc.

Some variants

Preaffectations



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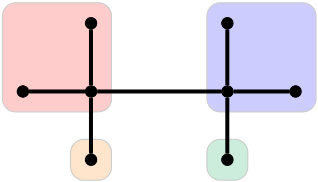

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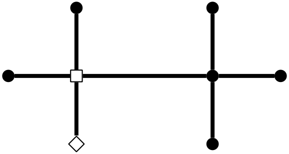

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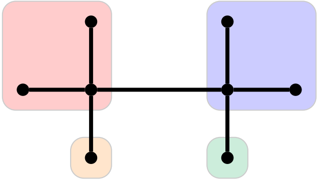

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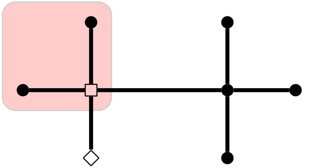

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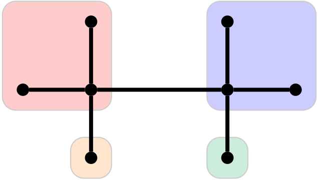

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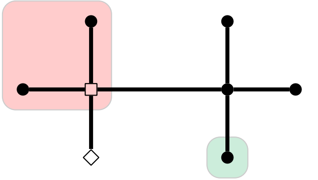

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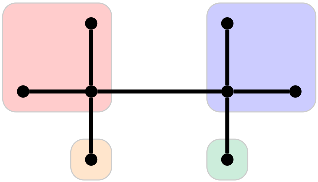

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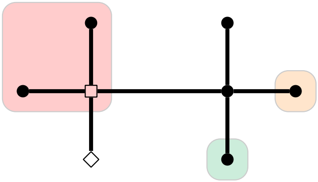

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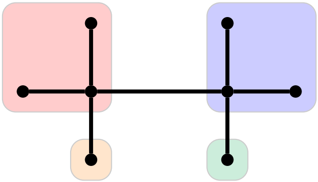

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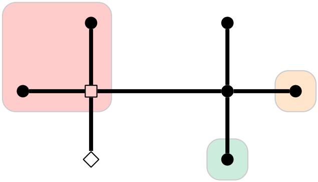

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Generalization of the latter results?

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Theorem [Baudon, B., Sopena, 2014]

For every $k \geq 1$ and $n \geq k$, there exist AP+k graphs on $\lceil \frac{n(k+1)}{2} \rceil$ edges.

Obtained (partly) by considering so-called **Harary graphs**.

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How to do that in general?

Other variants...

Still motivated by real-life modifications.

- **On-line AP:** Whenever a user requires resources, we provide them on-line. $\forall \lambda \in \{1, ..., V(G)\}$, there is a connected $G[S_\lambda]$ s.t. $G - S_\lambda$ is OL-AP.

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- ... and a few more algorithmic and structural things.

Perspectives, problems, etc.

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